

Extremal problems for convex lattice polytopes

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A sample problem

Jarník proved in 1926 that

if $\gamma \subset \mathbf{R}^2$ is a (closed) **strictly convex** curve of length ℓ , then

$$|\gamma \cap \mathbf{Z}^2| \leq \frac{3}{\sqrt[3]{2\pi}} \ell^{2/3} + O(\ell^{1/3}).$$

Here both the exponent $\frac{2}{3}$ and the constant $\frac{3}{\sqrt[3]{2\pi}}$ are best possible. Equivalently,

Theorem (Jarník 1926)

$$\lim_{\ell \rightarrow \infty} \max\{\ell^{-2/3} |\gamma \cap \mathbf{Z}^2| : \gamma \text{ is a convex..}\} = \frac{3}{\sqrt[3]{2\pi}}$$

convex lattice polygons appear instantly:

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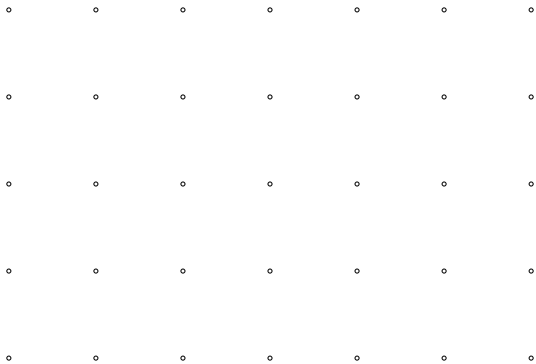
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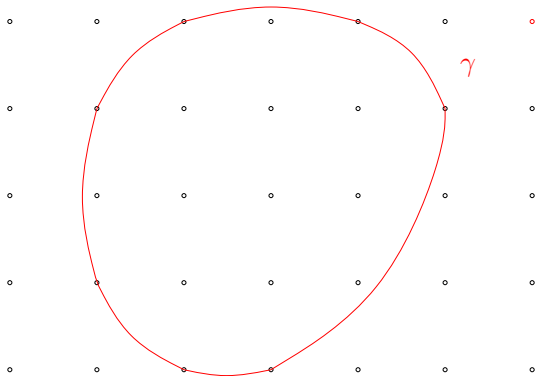
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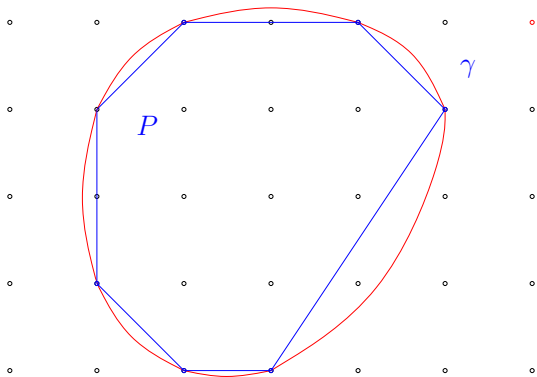
convex lattice polygons appear instantly:



The lattice \mathbf{Z}^2 (or \mathbf{Z}^d)



The strictly convex curve γ



The **convex lattice polygon** P whose vertex set is $\gamma \cap \mathbf{Z}^2$

In fact, $P = \text{conv}(\gamma \cap \mathbf{Z}^2)$. Jarník's result says that if P has $n = |\gamma \cap \mathbf{Z}^2|$ vertices, then

$$\ell > \text{per } P \geq \frac{\sqrt{6\pi}}{9} n^{3/2} + O(n^{3/4})$$

with best exponent $3/2$ and best constant $\frac{\sqrt{6\pi}}{9}$.

Theorem

With the min taken over all convex lattice polygons with n vertices

$$\lim_{n \rightarrow \infty} n^{-3/2} \min \text{per } P = \frac{\sqrt{6\pi}}{9}.$$

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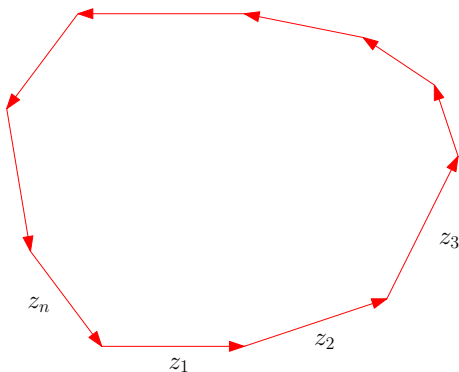
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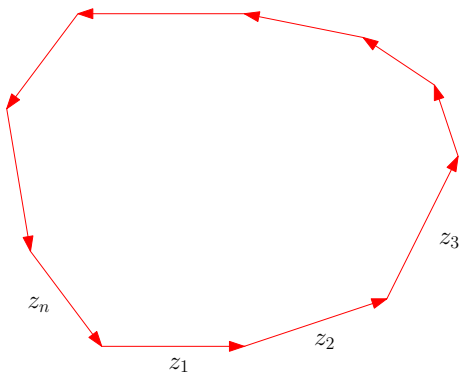
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P convex lattice n -gon with minimal perimeter, edges $z_1, z_2, \dots, z_n \in \mathbf{Z}^2$.



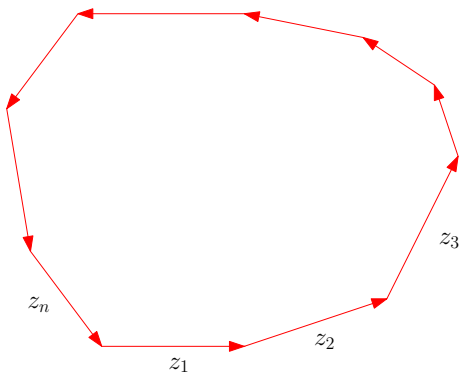
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- No z_i, z_j are parallel and same direction
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FACT: $z_1, \dots, z_n \in \mathbf{P}$ are distinct primitive vectors

Notation: $\mathbf{P} = \mathbf{P}^d \subset \mathbf{Z}^d$ set of primitive vectors

their density in \mathbf{Z}^2 is $6/\pi^2$

Let $U = \{u_1, \dots, u_n\}$ be the set of the n shortest primitive vectors.

$$\text{per } \mathbf{P} = \sum_1^n \|z_i\| \geq \sum_1^n \|u_i\|$$

$\sum_1^n \|u_i\|$ can be determined. With $r = \max \|u_i\|$

$$U \approx rB^2 \cap \mathbf{P} \text{ and } \frac{6}{\pi^2} r^2 \pi \approx n \text{ so } r \approx \sqrt{\frac{\pi n}{6}}.$$

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Similarly,

$$\begin{aligned} \text{per } P &\geq \sum_1^n \|u_i\| \approx \sum_{u \in rB^2 \cap \mathbf{P}} \|u\| \\ &\approx \frac{6}{\pi^2} \int_{rB^2} \|x\| dx \\ &\approx \frac{\sqrt{6\pi}}{9} n^{3/2}. \end{aligned}$$

Lower bound (for even n): choose the n shortest primitive vectors in pairs $-u, u$, so their sum is zero.

Order the vectors by increasing slope.

This gives the order of edges of a convex lattice polygon P and per $P \approx \frac{\sqrt{6\pi}}{9} n^{3/2}$.

For odd n ...

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REMARK. There is a **limit shape** of the minimizers (after scaling)

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$D \in \mathcal{K}^2$ with $0 \in D$ is the unit ball of a (non-symmetric) norm.
Let \mathcal{P} denote the family of all convex lattice polygons.
Each $P \in \mathcal{P}$ has a D -perimeter $\text{per}_D P$. Define

$$L_n(D) = \min\{\text{per}_D P : P \in \mathcal{P}, P \text{ has } n \text{ vertices}\}$$

Theorem (B.-Enriquez '10)

There is a convex set $P \subset \mathbf{R}^2$ such that the following holds. Let $P_n \in \mathcal{P}$ with n vertices be an arbitrary sequence of minimizers, of $L_n(D)$, translated so that their center of gravity is at the origin. Then the sequence $n^{-3/2}P_n$ tends to P .

P is unique

Proof: convex geometry, number theory, plus calculus of variation

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$\mathcal{P} = \mathcal{P}^d$ set of convex lattice polytopes,

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THE PROBLEMS

1. Minimal volume. Determine or estimate

$$V_d(n) = \min\{\text{Vol } P : P \in \mathcal{P}^d \text{ and } f_0(P) = n\}$$

2. Minimal surface area. Determine or estimate

$$S_d(n) = \min\{S(P) : P \in \mathcal{P}^d \text{ and } f_0(P) = n\}$$

just solved it for $d = 2$.

3. Minimal lattice width. Determine or estimate

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Definition

$K \in \mathcal{K}^d$, $z \in \mathbf{Z}^d$ and $z \neq 0$, then

$$w(K, z) = \max\{z \cdot (x - y) : x, y \in K\}.$$

The **lattice width** of K is

$$w(K) = \min\{w(K, z) : z \in \mathbf{Z}^d, z \neq 0\}.$$

How many parallel lattice hyperplanes meet K ?

FACT. For $P \in \mathcal{P}^d$, $w(P) + 1 =$ minimal number of parallel lattice lines meeting P .

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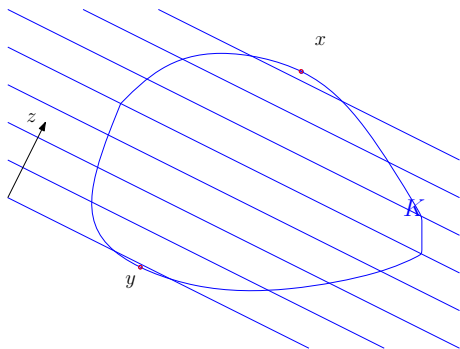
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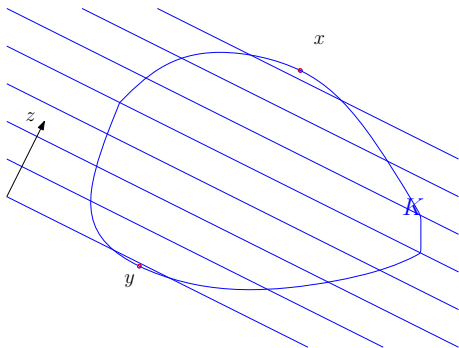
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$P, Q \in \mathcal{P}^d$ are equivalent if P can be carried to Q by a lattice preserving affine transformation.

Equivalent polytopes have the same volume.

Arnold's question. (1980) How many equivalence classes are there in \mathbf{R}^d , of volume $\leq V$?

not an extremal question yet ..

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$$\max\{f_0(P) : P \in \mathcal{P}^d, P \subset K\}.$$

equivalently, determine or estimate the maximal number of points in $K \cap \mathbf{Z}^d$ that are in convex position, i.e., none of them is in the convex hull of the others

answers: order of magnitude, asymptotic, precise..

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Theorem (Andrews '63)

If $P \in \mathcal{P}^d$ and $\text{Vol } P > 0$, then

$$f_0(P)^{\frac{d+1}{d-1}} \leq c_d \text{Vol } P.$$

or with better notation:

$$f_0(P)^{\frac{d+1}{d-1}} \ll \text{Vol } P.$$

Corollary

$$n^{\frac{d+1}{d-1}} \ll V_d(n).$$

Several proofs, none easy: Andrews '63, Arnold '80 ($d = 2$), Konyagin, Sevastyanov '84, ($d \geq 2$), W. Schmidt '86, B.-Vershik '92, B.-Larman '98, Reisner-Schütt-Werner '01, and more

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Several proofs, none easy: Andrews '63, Arnold '80 ($d = 2$), Konyagin, Sevastyanov '84, ($d \geq 2$), W. Schmidt '86, B.-Vershik '92, B.-Larman '98, Reisner-Schütt-Werner '01, and more

1. Minimal volume $V_d(n)$

Theorem (Andrews '63)

If $P \in \mathcal{P}^d$ and $\text{Vol } P > 0$, then

$$f_0(P)^{\frac{d+1}{d-1}} \leq c_d \text{Vol } P.$$

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Definition

A **tower** of $P \in \mathcal{P}^d$ is $F_0 \subset F_1 \subset \dots \subset F_{d-1}$ where F_i is an i -dim face of P . $T(P)$ = number of towers of P .

Theorem

If $P \in \mathcal{P}^d$ and $\text{Vol } P > 0$, then

$$T(P)^{\frac{d+1}{d-1}} \ll \text{Vol } P.$$

implies the same bound for $f_i(P)$.

OPEN PROBLEM. For all polytopes $P \in \mathcal{K}^d$

$$T(P) \ll f_0(P) + f_1(P) + \dots + f_{d-1}(P)????$$

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FACT. $n^{(d+1)/(d-1)}$ is best possible estimate

Example 1. (Arnold 80') G is the graph of the parabola $y = x^2$, $|x| \leq t$, and

$$P = P_t = \text{conv} (G \cap \mathbf{Z}^2).$$

Then $f_0(P) = 2t + 1$ and $\text{Area } P \approx \frac{2}{3}t^3$.

in d -dim, $G = G_t$ is given by $x_d = x_1^2 + \cdots + x_{d-1}^2 \leq t$,

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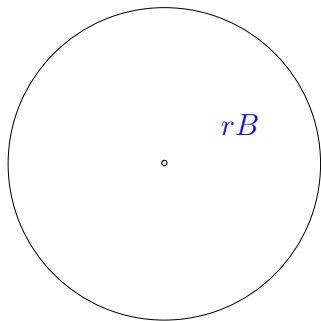
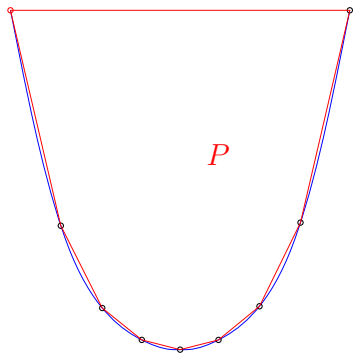
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the integer convex hull of rB^d

$\text{Vol } P_r \approx r^d$ implies via Andrews's theorem

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The **proof** uses the Flatness Theorem,
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If $P \subset B^d$ is a polytope with $f_0(P) \leq n$, then

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REMARK. Works for all $K \in \mathcal{K}^d$ (instead of B^d) with smooth enough boundary.

OPEN PROBLEM. Does $\lim n^{-\frac{d+1}{d-1}} V_d(n)$ exist????

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Isoperimetric inequality: For all $K \in \mathcal{K}^d$

$$\frac{S(K)^d}{(\text{Vol } K)^{d-1}} \geq \frac{S(B^d)^d}{(\text{Vol } B^d)^{d-1}}$$

implies $S(P) \gg (\text{Vol } P)^{(d-1)/d} \gg f_0(P)^{(d+1)/d}$

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$$n^{(d+1)/d} \ll S_d(n)$$

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3. Minimal lattice width $w_d(n)$

first $d = 2$. $w(P) + 1$ is the minimal number of consecutive lattice lines intersecting P .

each such line contains at most two vertices of $P \implies$
 $f_0(P) \leq 2(w(P) + 1)$

FACT. $w_2(n) = \lceil \frac{n}{2} \rceil - 1$

FACT. $w_d(n) = 1$

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4. Arnold's question

$P, Q \in \mathcal{P}^d$ are equivalent if a lattice preserving affine transformation maps P to Q .

FACT. $P \sim Q \implies f_0(P) = f_0(Q), w(P) = w(Q), \text{Vol } P = \text{Vol } Q.$

$N_d(V) =$ number of equivalent classes of $P \in \mathcal{P}^d$ with $\text{Vol } P \leq V$

$N_2(A)$ for $d = 2$

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Theorem (Arnold 1980)

$$A^{1/3} \ll \log N_2(A) \ll A^{1/3} \log A.$$

lower bound: let P be the polytope from Example 1 or 2.

Its vertex set $W \implies |W| \approx A^{1/3}$.

For each subset $U \subset W$, $\text{conv } U \in \mathcal{P}^2$.

there are $2^{|W|} \approx 2^{A^{1/3}}$ such subpolygons. Most of them distinct.

for the upper bound we need:

Lemma (Square lemma)

For every $P \in \mathbf{P}^2$ there is $Q \sim P$ which is contained in the square $[0, 36A]^2$.

So each equivalence class is represented in this square.

Proof follows from Andrews theorem + Square lemma

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So each equivalence class is represented in this square.

Proof follows from Andrews theorem + Square lemma

Theorem (Arnold 1980)

$$A^{1/3} \ll \log N_2(A) \ll A^{1/3} \log A.$$

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Its vertex set $W \implies |W| \approx A^{1/3}$.

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$$\text{Box}(\gamma) = \{x \in \mathbf{R}^d : 0 \leq x_i \leq \gamma_i, i = 1, \dots, d\}.$$

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Lemma (Key lemma)

The number of lattice polytopes contained in $\text{Box}(\gamma)$ is

$$\leq \exp \left(c_d (\text{Vol Box}(\gamma))^{\frac{d-1}{d+1}} \right)$$

ingredients:

- Minkowski's theorem: outer normals to the facets, of lengths equal to the surface area, determine P uniquely (up to translation),
- for a lattice polytope this outer normal vector is in $\frac{1}{(d-1)!} \mathbf{Z}^d$,
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FACT. The key lemma implies Andrews's theorem

Proof. $P \in \mathcal{P}^d$, $V = \text{Vol } P$, we assume $P \subset \text{Box}(\gamma)$ with $\text{Vol } \text{Box}(\gamma) \ll V$. Let $f_0(P) = n$.

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there are at least $2^n - 1$ distinct convex lattice polytopes in $\text{Box}(\gamma)$, the subpolytopes of P

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$$2^n - 1 \leq \exp \left(c_d (\text{Vol } \text{Box}(\gamma))^{\frac{d-1}{d+1}} \right)$$

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5. Maximal polytopes

better setting: $\mathbf{Z}_t = \frac{1}{t}\mathbf{Z}^d$ where t is large

$K \in \mathcal{K}^d$ is fixed with $\text{Vol } K = 1$, say

$\mathcal{P}(K, t)$ family of all convex \mathbf{Z}_t -lattice polytopes contained in K

$M(K, t) = \max\{f_0(P) : P \in \mathcal{P}(K, t)\}$,

same as maximal number of points in $\mathbf{Z}_t \cap K$ in convex position

Theorem

Suppose $K \in \mathcal{K}^d$ and $\text{Vol } K = 1$. Then

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Yes, when $d = 2$:

Theorem (B.-Prodromou '06)

When $K \subset \mathcal{K}^2$,

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$A(n) = \min\{\text{Area } P : P \in \mathcal{P}^d, f_0(P) = n\} = V_2(n)$ previous notation

Theorem (B.-Tokushige, '04)

$\lim n^{-3}A(n)$ exists and equals $0.0185067 \dots$ most likely.

FACT. $C \subset \mathbf{R}^2$ is an 0-symmetric convex body, and $|C \cap \mathbf{P}| = n$

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there is a unique (up to translation) convex lattice n -gon, $P(C)$, with edge set $C \cap \mathbf{P}$.

Proof: order the vectors in $C \cap \mathbf{P}$ by increasing slope...

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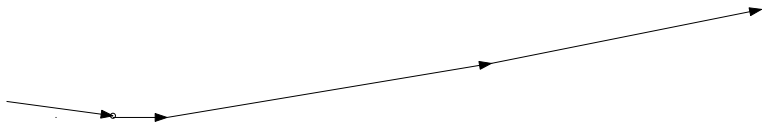
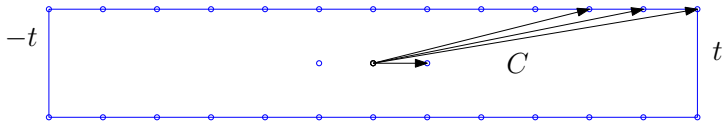
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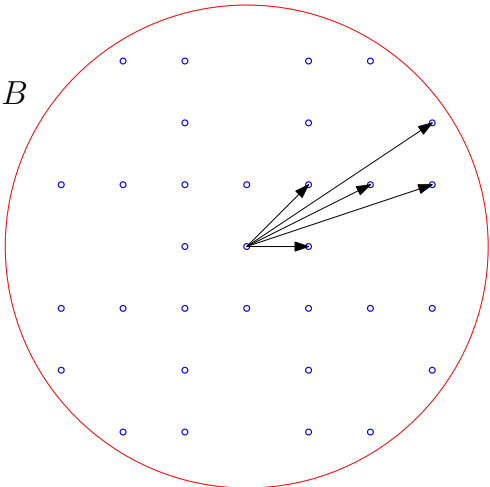
$$\begin{aligned} \text{Area } P(C) &= \left(\frac{1}{48} + o(1) \right) n^3 \\ &= (0.0204085 \dots + o(1)) n^3 \end{aligned}$$



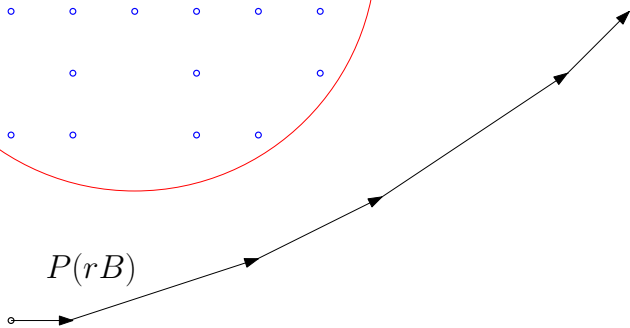
$C = rB^2$ with r chosen so that $|C \cap \mathbf{P}| = n \implies$

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rB



$P(rB)$



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Lemma (Reduction Lemma)

For even n , $A(n) = M(n)$.

Let C_n be a minimizer for $M(n)$, and w_n be its lattice width.

Theorem

There is a positive constant D such that

$$M(n) \geq \left(\frac{1}{54} - D \frac{\log w_n}{w_n} \right) n^3.$$

$D \approx 5000$

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or $w_n = \text{const}$ along a subsequence.

Determining $M(n)$ with side condition $w(C) = b$ leads to an
extremal problem $E(b)$ with b variables

which can be solved by a computer for fixed (not too large) b

$b = 8, 9, 14, 15, \dots$ gives $M(n)/n^3 < 1/54$

Enough to solve $E(b)$ for $b \leq 10^{10}$

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Arnold's problem in the plane:

$N(A)$ = number of equivalence classes in \mathcal{P}^2 of area $\leq A$.

OPEN PROBLEM. Does $\lim A^{-1/3} \log N(A)$ exist????